

# Chapter 10

## Finite State Machines Controller

### SKEE2263 Digital Systems

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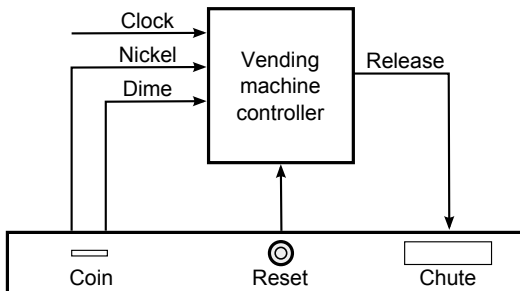
FEE, Universiti Teknologi Malaysia

April 27, 2016

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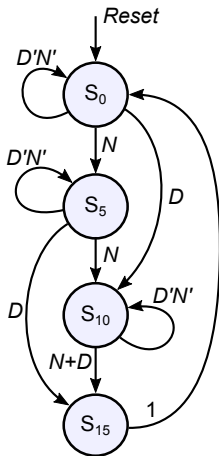
# Vending Machine Abstract View



## Vending Machine Specs

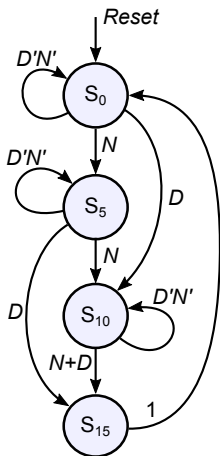
- Release gum after 15¢ is received
- Machine has single coin slot
- Accepts 5¢(nickel) and 10¢(dime), one at a time
- The controller's output causes gum to be dispensed

# Vending Machine Symbolic State Diagram



- Each state represents how much money has been deposited
- Variables:
  - $N$  = nickel, 10¢ detected
  - $D$  = dime, 5¢ detected
  - $R$  = release

# Vending Machine Symbolic State Table



Present State	Inputs		Next State	Output Release
	D	N		
0¢	0	0	0¢	0
	0	1	5¢	0
	1	0	10¢	0
	1	1	×	×
5¢	0	0	5¢	0
	0	1	10¢	0
	1	0	15¢	0
	1	1	×	×
10¢	0	0	10¢	0
	0	1	15¢	0
	1	0	15¢	0
	1	1	×	×
15¢	×	×	15¢	1

# Vending Machine Encoded State Table

State Encoding	
Symbolic	Encoded
0¢	00
5¢	01
10¢	10
15¢	11

Present State		Inputs		Next State	Output
$Q_1Q_0$		$D$	$N$	$D_1D_0$	$Release$
00		0	0	00	0
		0	1	01	0
		1	0	10	0
		1	1	××	×
01		0	0	01	0
		0	1	10	0
		1	0	11	0
		1	1	××	×
10		0	0	10	0
		0	1	11	0
		1	0	11	0
		1	1	××	×
11		×	×	00	1

## Vending Machine Equations

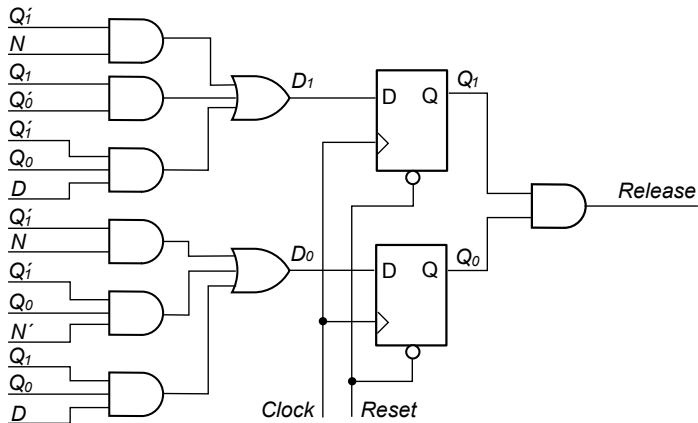
$$D_1 = Q_1' D + Q_1 Q_0' + Q_1' Q_0 N$$

$$D_0 = Q_1' N + Q_1' Q_0 N' + Q_1 Q_0 D$$

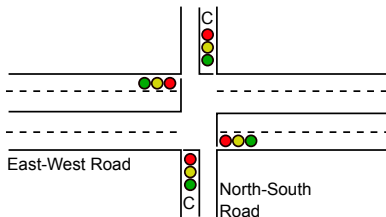
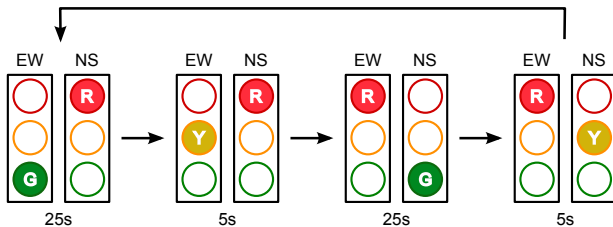
$$\textit{Release} = Q_1 Q_0$$



# Vending Machine Gate Level Schematic



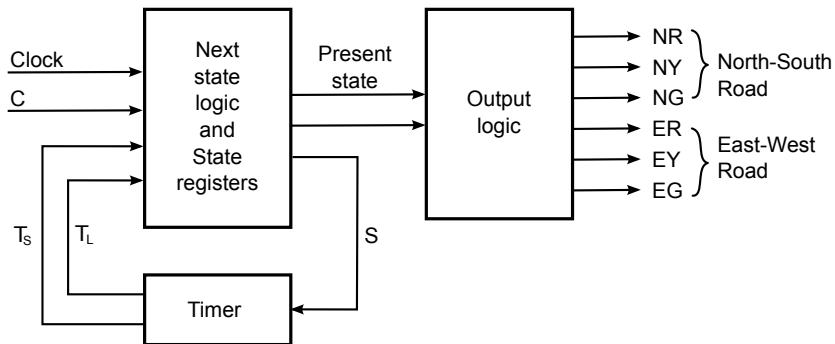
# Traffic Light Sequence



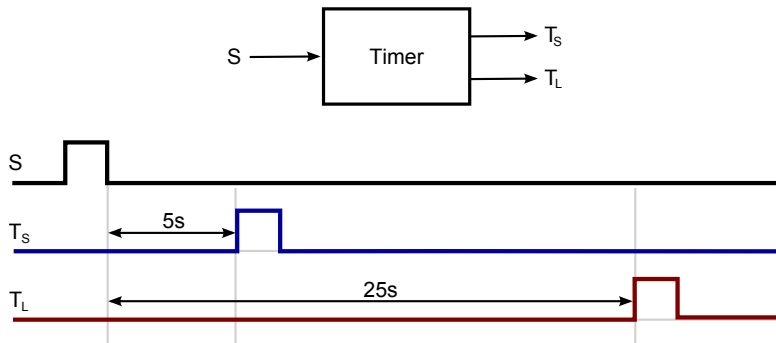
## Traffic Light Sequencing

State	East-West	North-South	Delay
EG	Green	Red	25 sec.
EY	Yellow	Red	5 sec.
NG	Red	Green	25 sec.
NY	Red	Yellow	5 sec.

# TLC Architecture

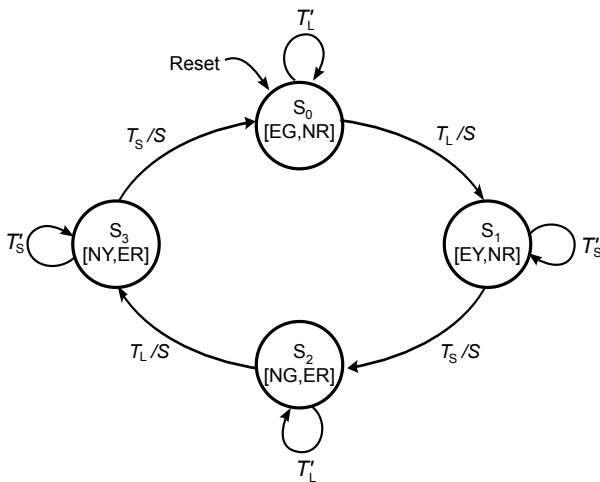


# Timer Module



- This is a separate state machine
- Just use as is. Don't worry about the internals.

# TLC Basic State Diagram

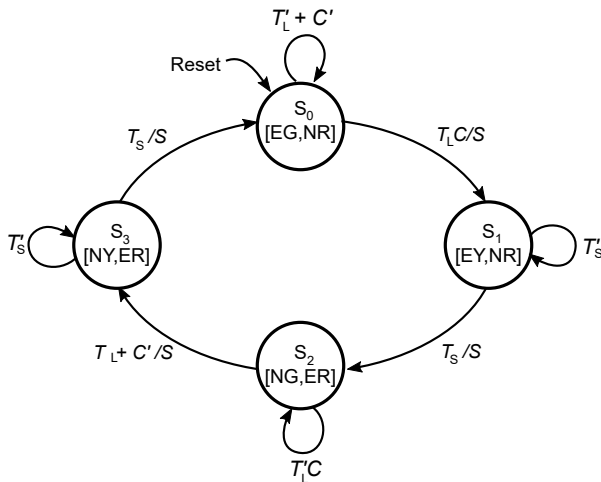


- Constant delay. Car detect (C) input not used.

# TLC Basic State Table

Present State	Inputs		Next State	Outputs							
	$Q_1Q_0$	$T_L$		$T_S$	$Q_1^+Q_0^+$	NR	NY	NG	ER	EY	EG
EG	0	X	EG	1	0	0	0	0	0	1	0
EG	1	X	EY	1	0	0	0	0	0	1	1
EY	X	0	EY	1	0	0	0	0	1	0	0
EY	X	1	NG	1	0	0	0	0	1	0	1
NG	0	X	NG	0	0	1	1	1	0	0	0
NG	1	X	NY	0	0	1	1	1	0	0	1
NY	X	0	NY	0	1	0	1	1	0	0	0
NY	X	1	EG	0	1	0	1	1	0	0	1

# TLC with Car Detection



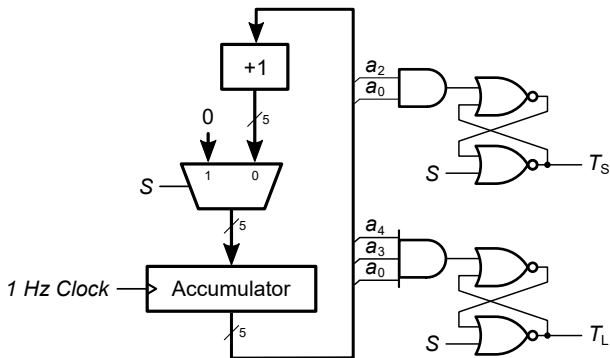
- If  $C = 0$ , wait in state EG.



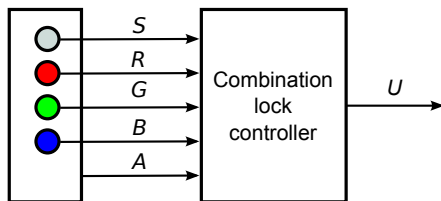
# TLC Car Detection State Table

Present State	Inputs			Next State	Outputs							
	$Q_1Q_0$	$C$	$T_L$		$T_S$	$Q_1^+Q_0^+$	NR	NY	NG	ER	EY	EG
EG	0	X	X	EG	1	0	0	0	0	0	1	0
EG	X	0	X	EG	1	0	0	0	0	0	1	0
EG	1	1	X	EY	1	0	0	0	0	0	1	1
EY	X	X	0	EY	1	0	0	0	0	1	0	0
EY	X	X	1	NG	1	0	0	0	0	1	0	1
NG	1	0	X	NG	0	0	1	1	1	0	0	0
NG	0	X	X	NY	0	0	1	1	1	0	0	0
NG	X	1	X	NY	0	0	1	1	1	0	0	1
NY	X	X	0	NY	0	1	0	1	1	0	0	0
NY	X	X	1	EG	0	1	0	1	1	0	0	1

# TLC Timer Design



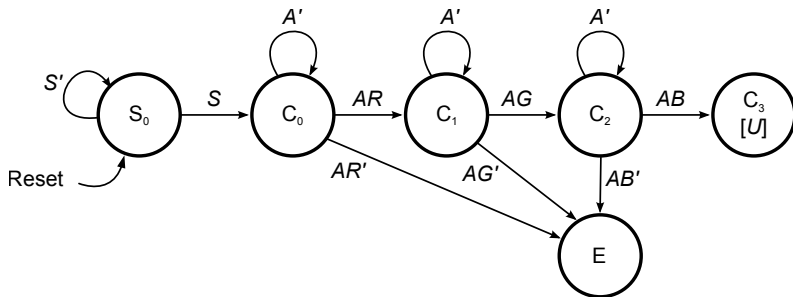
## Combination Lock Top Level



- The controller unlocks the door when buttons are pressed in this sequence:

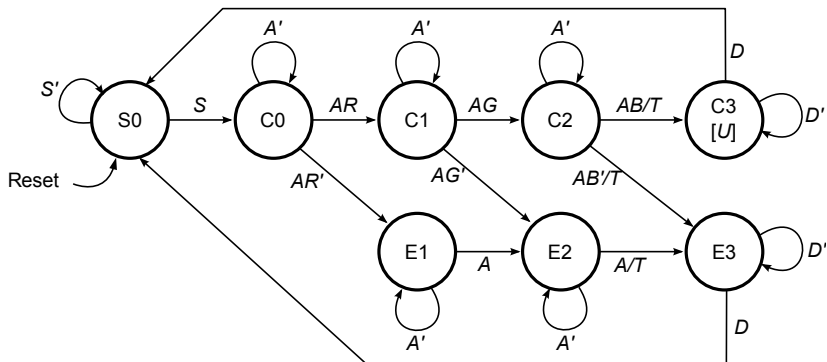
Start → Red → Green → Blue

## Combination Lock 1st Try State Diagram



- It works but lock is easily breakable
- Why?

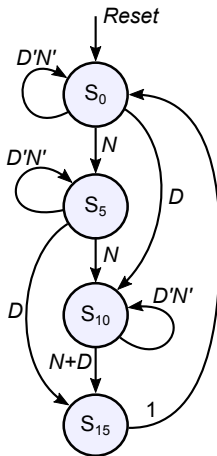
# A More Robust Combination Lock



# One-Hot Encoding

- One state = one flip-flop
  - example: 8 states = 3 FF for straight binary or Gray encoding
  - 8 states = 8 FF for one hot encoding
- Only 1 FF high for any given state, all other low
- The high FF is called the hot FF
- Wastes FF, but on FPGAs FF are plenty
- **Simpler to generate next state logic**

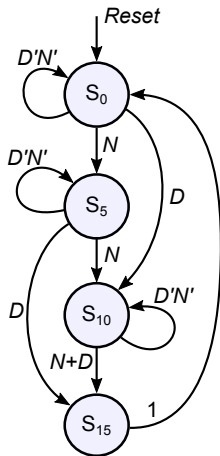
# One-Hot Moore Vending Machine



One-Hot State Assignment

Symbolic	One-Hot Encoding
$S_0 = 0\text{¢}$	1000
$S_5 = 5\text{¢}$	0100
$S_{10} = 10\text{¢}$	0010
$S_{15} = 15\text{¢}$	0001

# One-Hot Moore Vending Machine



Next state equations:

$$S_0^+ = S_0 D' N' + S_{15}$$

$$S_5^+ = S_0 N + S_5 D' N'$$

$$S_{10}^+ = S_0 D + S_5 N + S_{10} D' N'$$

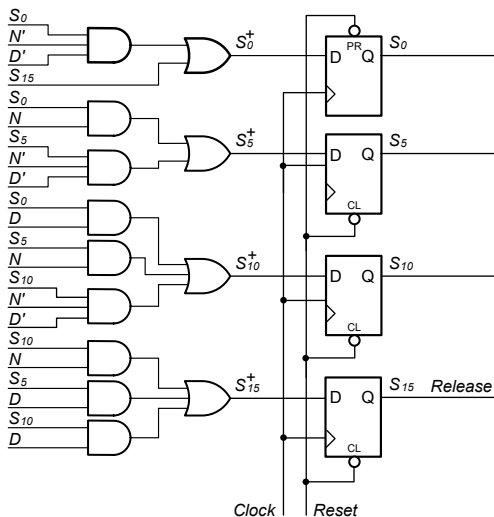
$$S_{15}^+ = S_5 D + S_{10} N + S_{10} D$$

Output equation is:

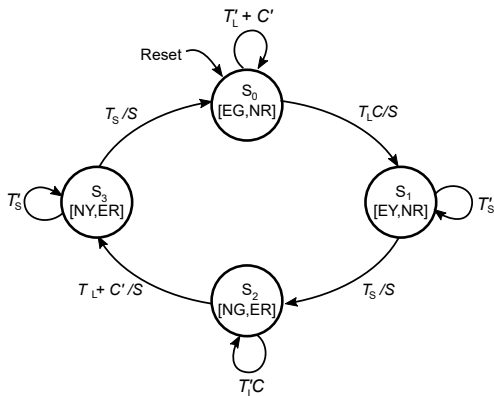
$$Release = S_{15}$$



# One-Hot Moore Vending Machine



# One-Hot Traffic Light Controller



### One-Hot State Assignment

Symbolic	Encoding
$S_0 = 0\phi$	1000
$S_1 = 5\phi$	0100
$S_2 = 10\phi$	0010
$S_3 = 15\phi$	0001

# One-Hot Traffic Light Controller

Next state equations:

$$S_0^+ = S_0(T'_L + C') + S_3T_S = S_0T'_L + S_0C' + S_3T_S$$

$$S_1^+ = S_0T_L C + S_1T'_S$$

$$S_2^+ = S_1T_S + S_2T'_L.C$$

$$S_3^+ = S_2(T_L + C') + S_3T'_S = S_2T_L + S_2C' + S_3T'_S$$

Output equations:

$$S = S_0T_L C + S_1T_S + S_2(T_L + C') + S_3T_S$$

$$EG = S_0$$

$$EY = S_1$$

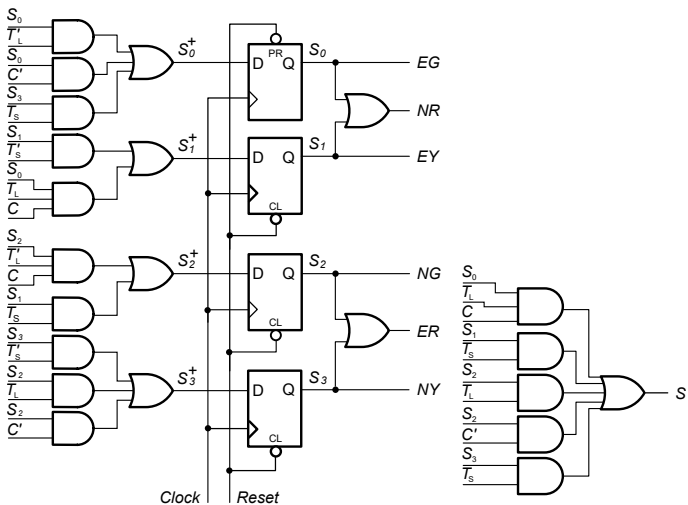
$$ER = S_2 + S_3$$

$$NG = S_2$$

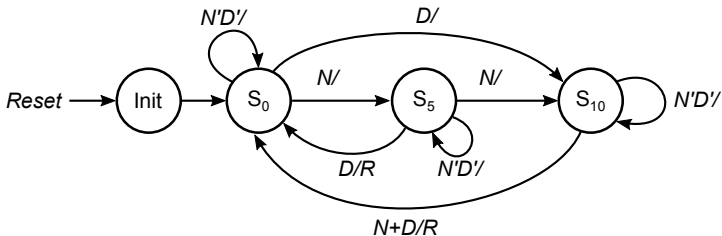
$$NY = S_3$$

$$NR = S_0 + S_1$$

# One-Hot Traffic Light Controller



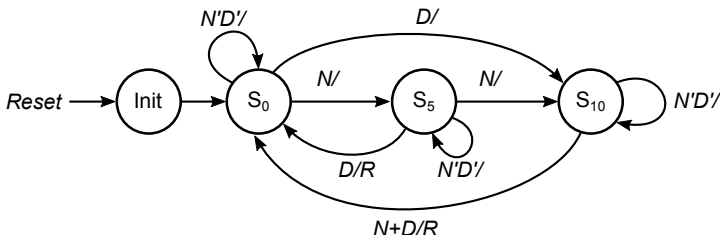
# Almost One-Hot Mealy-type Vending Machine



One-Hot State Assignment

Symbolic	One-Hot Encoding
$Init =$ Initialization	000
$S_0 =$ 0¢	100
$S_5 =$ 5¢	010
$S_{10} =$ 10¢	001

# Almost One-Hot Mealy-type Vending Machine



Next state & output equations:

$$S_0 = Init + S_0 N' D' + S_5 D + S_{10} N + S_{10} D$$

$$S_5 = S_0 N + S_5 N' D'$$

$$S_{10} = S_0 D + S_5 N + S_{10} N' D'$$

$$R = S_5 D + S_{10} D + S_{10} N$$

There is no next state equation for *Init*. Why?

# Almost One-Hot Mealy-type Vending Machine

